Eliminating Exception Constraints of Java Programs for IA-64



<http://www.trl.ibm.com/projects/jit/index_e.htm>

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Goal of the Paper

Motivation

- Enable to perform code motion to exploit *instruction level* parallelism (ILP) of IA-64 for Java
- Enable to perform only beneficial speculative code motion
- Our approach "exception speculation" using speculative code motion
 - ► Perform exception speculation on *directed acyclic graph* (DAG)
- Experimental results
- Summary



A Running ExampleJava program and bytecode

iconst_1

ireturn

iadd

```
Java program
int foo(int a[], int i) {
   return a[i] + 1;
}
Bytecode
iaload
   may throw a Java exception.
```



Intermediate Representation Java language introduces many exception checks





Problems in Java

An exception dependence between exception check and load suppresses code motion



Control Speculation in IA-64
 An speculative load instruction allows dependant loads to issue a load earlier before the conditional branch is resolved.



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Our Approach - Exception SpeculationEliminate exception dependence edges from each load



Why We Distinguish Between Control and Exception Speculation

- Reduce the size of IR by not splitting basic blocks
 - ► We do not handle exception dependence as control dependence.
 - In our experiments, # of basic blocks can be increased by a factor of four without using exception dependence edges.
- Estimate the benefit of exception speculation along the exception dependence edge.
 - The code can be moved speculatively only when it is beneficial on the DAG.



Where We Perform Exception Speculation





Algorithm Outline

1. Decide whether a load can be moved speculatively

When Delay(n) is set only by exception dependence, where n is an instruction.

 $Delay(n) = \max_{m \in Pred(n, DAG)} Delay(m) + Latency(m)$

2. Determine a speculative chain

- Load and the succeeding instructions w/o side effect
- 3. Eliminate and connect exception dependence edges
 - ► Restructure a DAG to issue a load earlier.
- 4. Create dependence edges
 - Maintain edges to preserve the correctness



Our DAG for Exception SpeculationBefore eliminating exception dependence edge_arg1

Java Program
int foo(int a[], int i) {
 return a[i] + 1;
}

Intermediate Representation

nullcheck	a
len	v1 = [a]
add	v2 = a, 16
shiftl	v3 = i, 2
boundcheck	i < v1
boundcheck ld	i < v1 v4 = [v2+v3]
boundcheck ld add	i < v1 v4 = [v2+v3] v5 = v4, 1









Decision to Perform Exception Speculation

- Calculate the maximum possible delay to execute *load*
 - Perform exception speculation if the time set by exception dependence is the slowest.





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Determine Speculative Chain

Determine a chain of instructions that have no side effects as a speculative chain







Dependence Edges



Experimental Results

- Measurements for:
 - Performance improvement
 - ► Code size expansion
- Benchmarks
 - ► Java Grande Benchmark Version 2.0
 - ► SPECjvm98
- Environments
 - ► IBM Developers Kit for IA-64, Java Technology Edition, 1.3
 - ► 2-way 800MHz Itanium with 2GB memory
 - Windows XP Advanced Server









Propose a new solution "exception speculation"

- Eliminate constraint of a load instruction by exception dependences on a DAG representation.
- Perform speculative code motion based on cost-benefit analysis.
- Show the effectiveness using a production Java JIT compiler



Thanks !!

Let's take a coffee break.



